
Software Model Checking Improving Security of a Billion Computers

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Acknowledgments

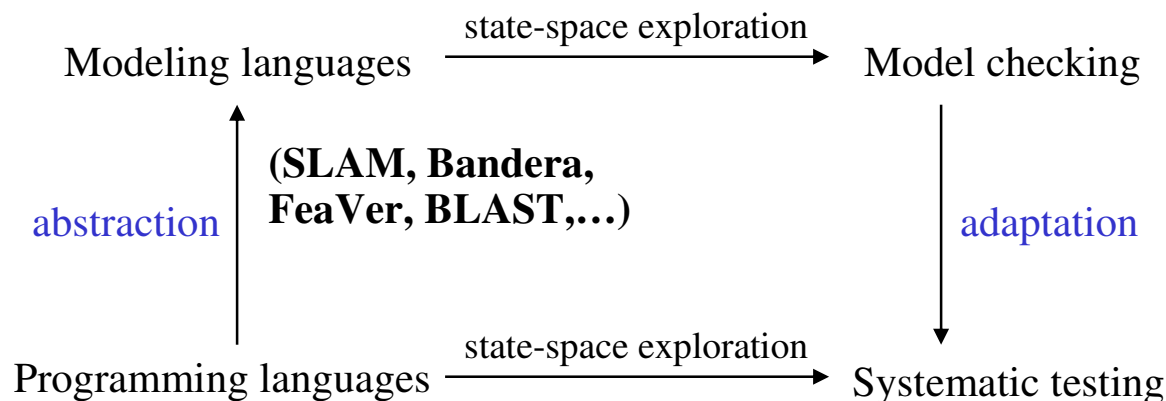
- Joint work with **Michael Levin** (CSE) and others:
 - Chris Marsh, Lei Fang, Stuart de Jong (CSE)
 - interns Dennis Jeffries (06), David Molnar (07), Adam Kiezun (07), Bassem Elkarablieh (08), ...
- Thanks to the entire SAGE team and users !
 - MSR: Ella Bounimova,...
 - Z3: Nikolaj Bjorner, Leonardo de Moura,...
 - WEX (Windows): Nick Bartmon, Eric Douglas,...
 - Office: Tom Gallagher, Octavian Timofte,...
 - SAGE users all across Microsoft!

References

- see <http://research.microsoft.com/users/pg>
 - DART: Directed Automated Random Testing, with N. Klarlund and K. Sen, PLDI'2005
 - Compositional Dynamic Test Generation, POPL'2007
 - Automated Whitebox Fuzz Testing, with M. Levin and D. Molnar, NDSS'2008
 - Demand-Driven Compositional Symbolic Execution, with S. Anand and N. Tillmann, TACAS'2008
 - Grammar-Based Whitebox Fuzzing, with A. Kiezun and M. Levin, PLDI'2008
 - Active Property Checking, with M. Levin and D. Molnar, EMSOFT'2008
 - Precise Pointer Reasoning for Dynamic Test Generation, with B. Elkarablieh and M. Levin, ISSTA'2009

A Brief History of Software Model Checking

- How to apply model checking to analyze **software**?
 - “Real” programming languages (e.g., C, C++, Java),
 - “Real” size (e.g., 100,000's lines of code).
- Two main approaches to software model checking:



Concurrency: **VeriSoft, JPF, CMC, Bogor, CHES**,...
Data inputs: **DART**, **EXE**, **SAGE**,...

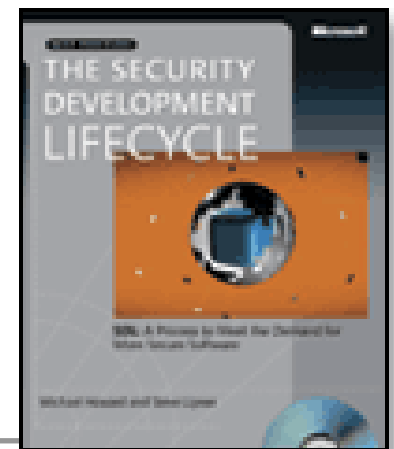
Killer app: security biggest impact to date!

Security is Critical (to Microsoft)

- Software security bugs can be very expensive:
 - Cost of each Microsoft Security Bulletin: \$Millions
 - Cost due to worms (Slammer, CodeRed, Blaster, etc.): \$Billions
- Most security exploits are initiated via files or packets
 - Ex: Internet Explorer parses dozens of file formats
- Security testing: "hunting for million-dollar bugs"
 - Write A/V (always exploitable), Read A/V (sometimes exploitable), NULL-pointer dereference, division-by-zero (harder to exploit but still DOS attacks), etc.

Hunting for Security Bugs

- Main techniques used by “black hats”:
 - Code inspection (of binaries) and
 - Blackbox fuzz testing
- Blackbox fuzz testing:
 - A form of blackbox random testing [Miller+90]
 - Randomly **fuzz** (=modify) a well-formed input
 - Grammar-based fuzzing: rules that encode “well-formed”ness + heuristics about how to fuzz (e.g., using probabilistic weights)
- **Heavily** used in security testing
 - Ex: July 2006 “Month of Browser Bugs”
 - Simple yet effective: many bugs found this way...
 - At Microsoft, fuzzing is mandated by the SDL



Blackbox Fuzzing

- Examples: Peach, Protos, Spike, Autodafe, etc.
- Why so many blackbox fuzzers?
 - Because anyone can write (a simple) one in a week-end!
 - Conceptually simple, yet effective...
- Sophistication is in the “add-on”
 - Test harnesses (e.g., for packet fuzzing)
 - Grammars (for specific input formats)
- Note: usually, no principled “spec-based” test generation
 - No attempt to cover each state/rule in the grammar
 - When probabilities, no global optimization (simply random walks)

Introducing Whitebox Fuzzing

- Idea: mix fuzz testing with dynamic test generation
 - Symbolic execution
 - Collect constraints on inputs
 - Negate those, solve with constraint solver, generate new inputs
 - do "systematic dynamic test generation" (=DART)
- Whitebox Fuzzing = "DART meets Fuzz"
Two Parts:
 1. Foundation: DART (Directed Automated Random Testing)
 2. Key extensions ("Whitebox Fuzzing"), implemented in SAGE

Automatic Code-Driven Test Generation

Problem:

Given a sequential program with a set of input parameters,
generate a set of inputs that maximizes code coverage

= "automate test generation using program analysis"

This is **not** "model-based testing"
(= generate tests from an FSM spec)

How? (1) Static Test Generation

- Static analysis to partition the program's input space [King76,...]
- Ineffective whenever symbolic reasoning is not possible
 - which is frequent in practice... (pointer manipulations, complex arithmetic, calls to complex OS or library functions, etc.)

Example:

```
int obscure(int x, int y) {  
    if (x==hash(y)) error();  
    return 0;  
}
```

Can't statically generate
values for x and y
that satisfy "x==hash(y)" !

How? (2) **Dynamic** Test Generation

- Run the program (starting with some random inputs), gather constraints on inputs at conditional statements, use a constraint solver to generate new test inputs
- Repeat until a specific program statement is reached [Korel90,...]
- Or repeat to try to cover **ALL** feasible program paths:
DART = Directed Automated Random Testing
= systematic dynamic test generation [PLDI'05,...]
 - detect crashes, assertion violations, use runtime checkers (Purify,...)

DART = Directed Automated Random Testing

Example:

```
int obscure(int x, int y) {  
    if (x==hash(y)) error();  
    return 0;  
}
```

Run 1 :- start with (random) x=33, y=42

- execute concretely and symbolically:

if (33 != 567) | if (x != hash(y))

constraint too complex

simplify it: x != 567

- solve: x==567 solution: x=567

- new test input: x=567, y=42

Run 2 : the other branch is executed

All program paths are now covered !

- Observations:

- Dynamic test generation extends static test generation with additional runtime information: it is more powerful
- The number of program paths can be infinite: may not terminate!
- Still, DART works well for small programs (1,000s LOC)
- Significantly improves code coverage vs. random testing

DART Implementations

- Defined by symbolic execution, constraint generation and solving
 - Languages: C, Java, x86, .NET,...
 - Theories: linear arith., bit-vectors, arrays, uninterpreted functions,...
 - Solvers: lp_solve, CVCLite, STP, Disolver, Z3,...
- Examples of tools/systems implementing DART:
 - EXE/EGT (Stanford): independent ['05-'06] closely related work
 - CUTE = same as first DART implementation done at Bell Labs
 - SAGE (CSE/MSR) for x86 binaries and merges it with "fuzz" testing for finding security bugs (more later)
 - PEX (MSR) for .NET binaries in conjunction with "parameterized-unit tests" for unit testing of .NET programs
 - YOGI (MSR) for checking the feasibility of program paths generated statically using a SLAM-like tool
 - Vigilante (MSR) for generating worm filters
 - BitScope (CMU/Berkeley) for malware analysis
 - CatchConv (Berkeley) focus on integer overflows
 - Splat (UCLA) focus on fast detection of buffer overflows
 - Apollo (MIT) for testing web applications

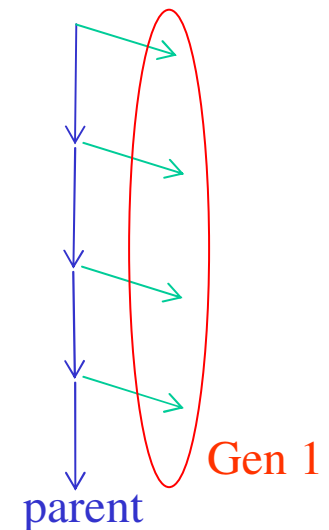
...and more!

DART Summary

- DART attempts to exercise all paths (like model checking)
 - Covering a single specific assertion (verification): hard problem (often intractable)
 - Maximize path coverage while checking thousands of assertions all over: easier problem (optimization, best-effort, tractable)
 - Better coverage than pure random testing (with directed search)
- DART can work around limitations of symbolic execution
 - Symbolic execution is an adjunct to concrete execution
 - Concrete values are used to simplify unmanageable symbolic expressions
 - Randomization helps where automated reasoning is difficult
- Comparison with static analysis:
 - No false alarms (more precise) but may not terminate (less coverage)
 - "Dualizes" static analysis: static **may** vs. DART **must**
 - Whenever symbolic exec is too hard, under-approx with concrete values
 - If symbolic execution is perfect, no approx needed: both coincide!

Whitebox Fuzzing [NDSS'08]

- Whitebox Fuzzing = "DART meets Fuzz"
- Apply DART to large applications (not unit)
- Start with a well-formed input (not random)
- Combine with a **generational** search (not DFS)
 - Negate 1-by-1 **each** constraint in a path constraint
 - Generate **many** children for each parent run
 - Challenge **all** the layers of the application sooner
 - Leverage expensive symbolic execution
- Search spaces are **huge**, the search is **partial**... yet **effective** at finding bugs !



Example

```
void top(char input[4])
```

```
{
```

```
    int cnt = 0;
```

```
    if (input[0] == 'b') cnt++;
```

```
    if (input[1] == 'a') cnt++;
```

```
    if (input[2] == 'd') cnt++;
```

```
    if (input[3] == '!') cnt++;
```

```
    if (cnt >= 3) crash();
```

```
}
```

input = "good"

Path constraint:

$I_0 \neq \text{'b'}$

$I_1 \neq \text{'a'}$

$I_2 \neq \text{'d'}$

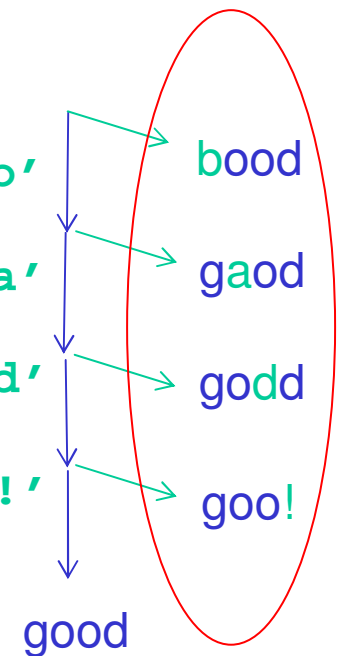
$I_3 \neq \text{'!'}$

$I_0 = \text{'b'}$

$I_1 = \text{'a'}$

$I_2 = \text{'d'}$

$I_3 = \text{'!'}$



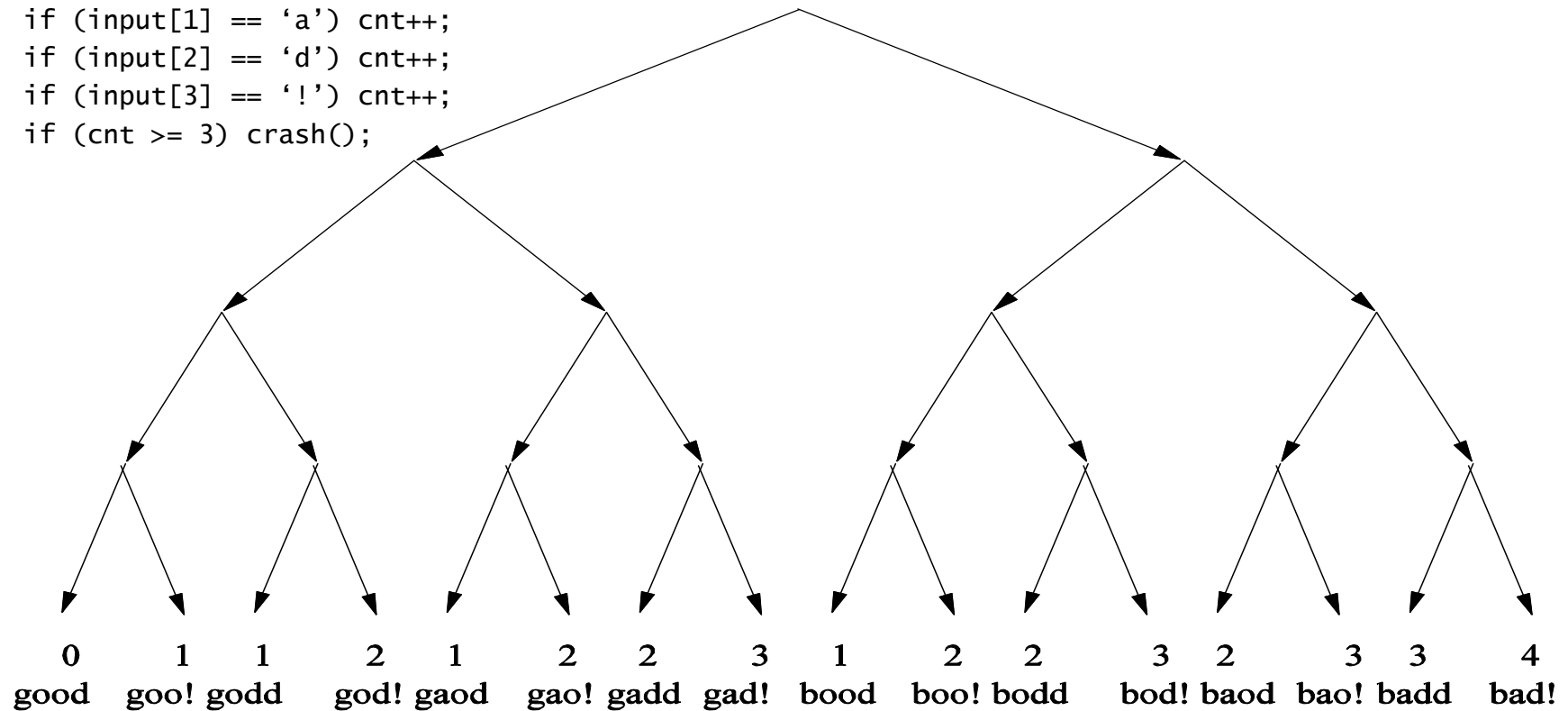
Gen 1

Negate each constraint in path constraint

Solve new constraint new input

The Search Space

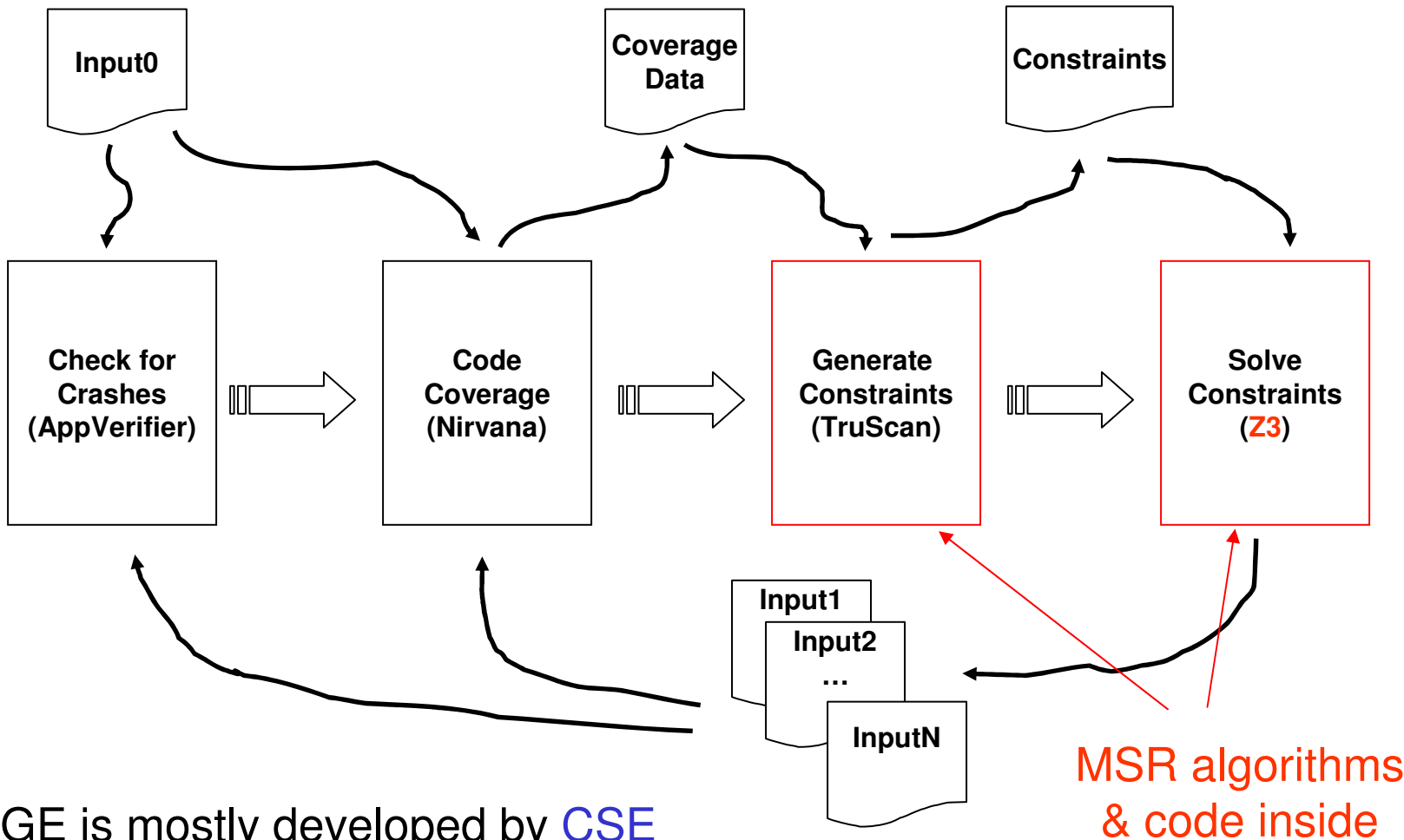
```
void top(char input[4])
{
    int cnt = 0;
    if (input[0] == 'b') cnt++;
    if (input[1] == 'a') cnt++;
    if (input[2] == 'd') cnt++;
    if (input[3] == '!') cnt++;
    if (cnt >= 3) crash();
}
```



SAGE (Scalable Automated Guided Execution)

- Generational search introduced in SAGE
- Performs symbolic execution of x86 execution traces
 - Builds on Nirvana, iDNA and TruScan for x86 analysis
 - Don't care about language or build process
 - Easy to test new applications, no interference possible
- Can analyse any file-reading Windows applications
- Several optimizations to handle huge execution traces
 - Constraint caching and common subexpression elimination
 - Unrelated constraint optimization
 - Constraint subsumption for constraints from input-bound loops
 - "Flip-count" limit (to prevent endless loop expansions)

SAGE Architecture



Some Experiments

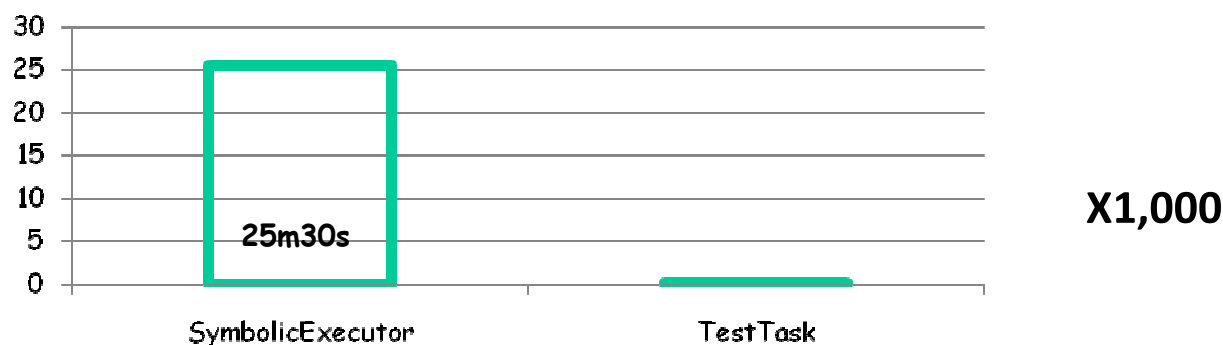
Most much (100x) bigger than ever tried before!

- Seven applications - 10 hours search each

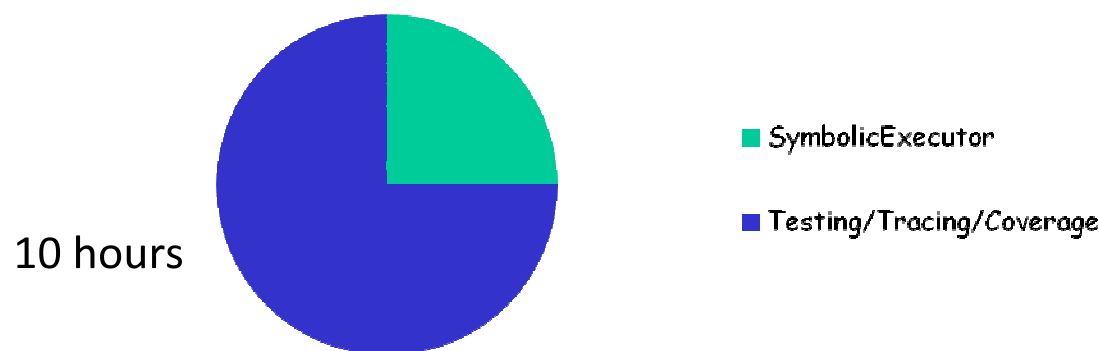
| App Tested | #Tests | Mean Depth | Mean #Instr. | Mean Input Size |
|------------------------|--------|------------|--------------|-----------------|
| ANI | 11468 | 178 | 2,066,087 | 5,400 |
| Media1 | 6890 | 73 | 3,409,376 | 65,536 |
| Media2 | 1045 | 1100 | 271,432,489 | 27,335 |
| Media3 | 2266 | 608 | 54,644,652 | 30,833 |
| Media4 | 909 | 883 | 133,685,240 | 22,209 |
| Compressed File Format | 1527 | 65 | 480,435 | 634 |
| OfficeApp | 3008 | 6502 | 923,731,248 | 45,064 |

Generational Search Leverages Symbolic Execution

- Each symbolic execution is expensive



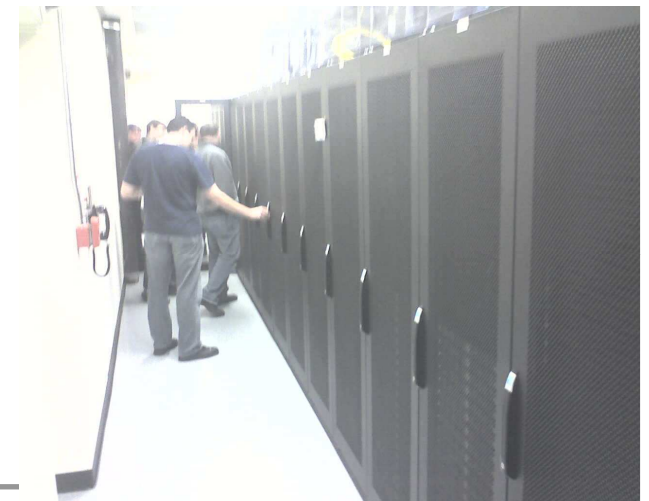
- Yet, symbolic execution does not dominate search time



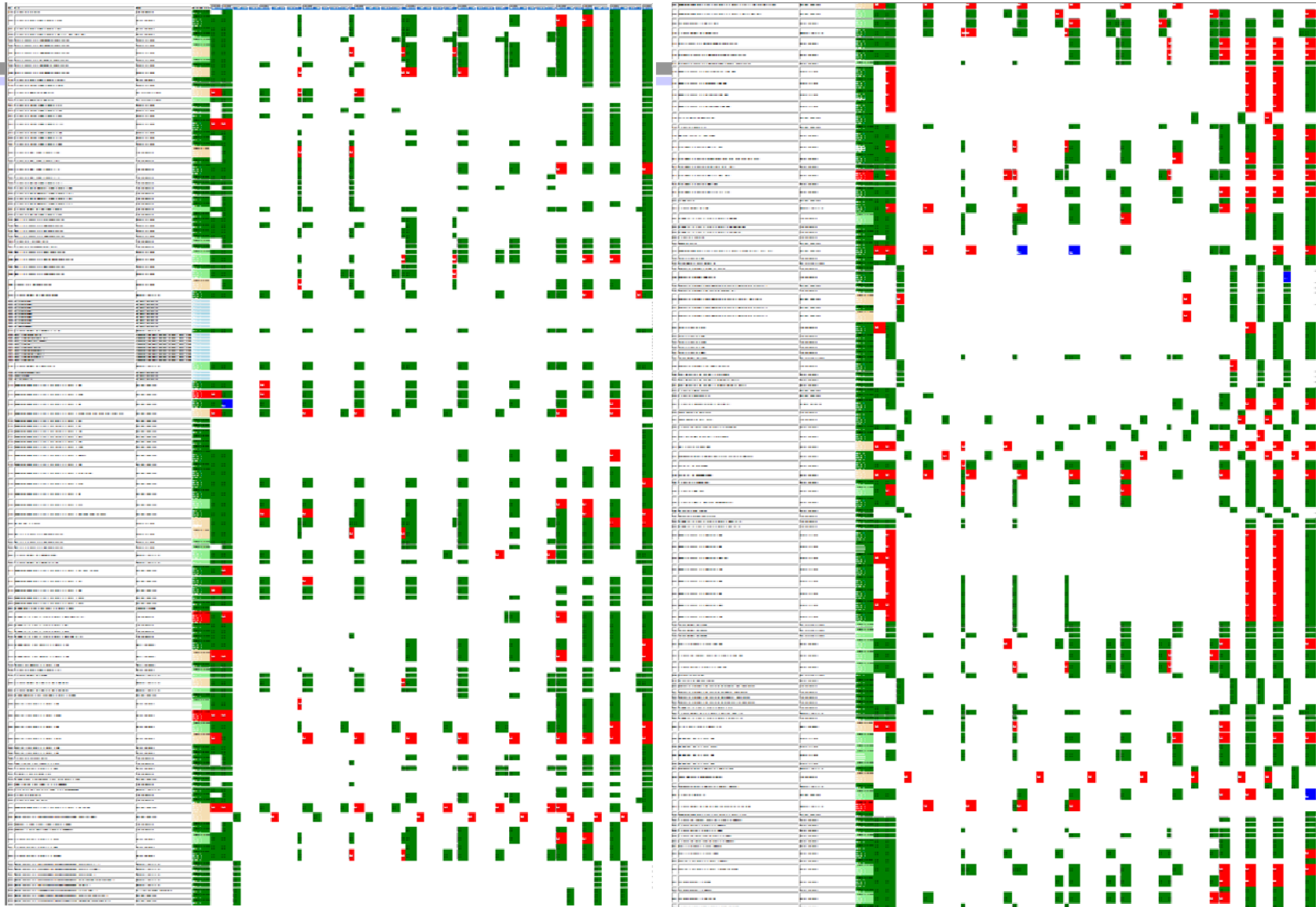
SAGE Results

Since April'07 1st release: many new security bugs found
(missed by blackbox fuzzers, static analysis)

- Apps: image processors, media players, file decoders,...
- Bugs: Write A/Vs, Read A/Vs, Crashes,...
- Many triaged as "security critical, severity 1, priority 1"
(would trigger Microsoft security bulletin if known outside MS)
- Example: WEX Security team for Win7
 - Dedicated fuzzing lab with 100s machines
 - 100s apps (deployed on 1billion+ computers)
 - ~1/3 of **all** fuzzing bugs found by SAGE !
- SAGE = **gold** medal at Fuzzing Olympics
organized by SWI at BlueHat'08 (Oct'08)
- Credit due to entire SAGE team + users !

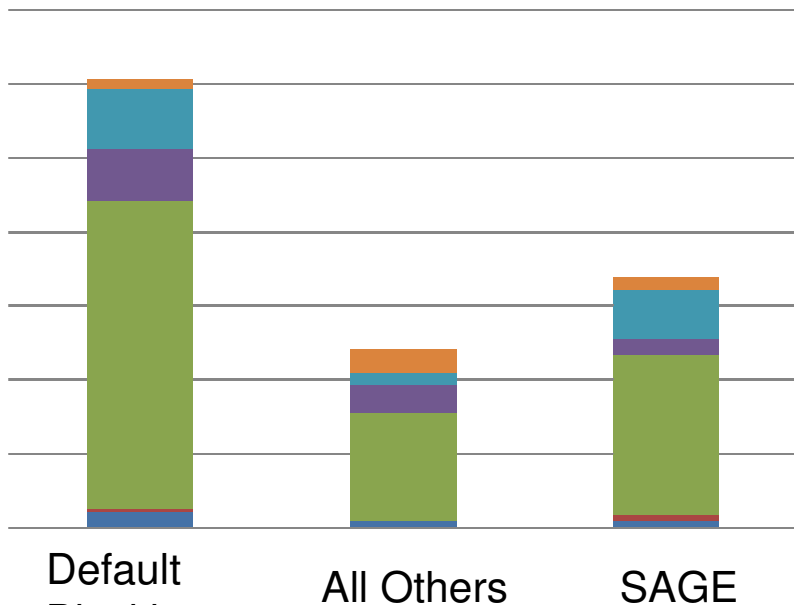


WEX Fuzz Dashboard Snippet



WEX Fuzzing Lab Bug Yield for Win7

How fuzzing bugs found (2006-2009) :



Default
Blackbox
Fuzzer
+ Regression

All Others

SAGE

SAGE is running 24/7 on 100s machines:
“the largest usage ever of any SMT solver”
N. Bjorner + L. de Moura (MSR, Z3 authors)

- 100s of apps, total number of fuzzing bugs is confidential
- But SAGE didn't exist in 2006
- Since 2007 (SAGE 1st release),
~1/3 bugs found by SAGE
- But SAGE currently deployed on only ~2/3 of those apps
- Normalizing the data by 2/3,
SAGE found ~1/2 bugs
- SAGE is more CPU expensive, so it is run last in the lab, so all SAGE bugs were missed by everything else!

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- *SAGE* generates a crashing test for *Media1* parser:

```
00000000h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000060h: 00 00 00 00 ; ....
```

Generation 0 – seed file

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 00 00 00 00 00 00 00 00 00 00 00 00 ; RIFF.....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 1

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF... *** .....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; .....
```

Generation 2

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF[REDACTED]...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 3

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 00 00 00 ; ....strh.....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 4

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 5

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 00 00 00 00 ; ....strf.....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 6

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 28 00 00 00 ; ....strf....(
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 7

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 28 00 00 00 ; ....strf....(...)
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 C9 9D E4 4E ; .....E\an
00000060h: 00 00 00 00 ; ....
```

Generation 8

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 28 00 00 00 ; ....strf....(...)
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 9

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser:

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 B2 75 76 3A 28 00 00 00 ; ....strf2uv:(...
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 10 – crash bucket 1212954973!

Found after only 3 generations starting from seed3 file on next slide

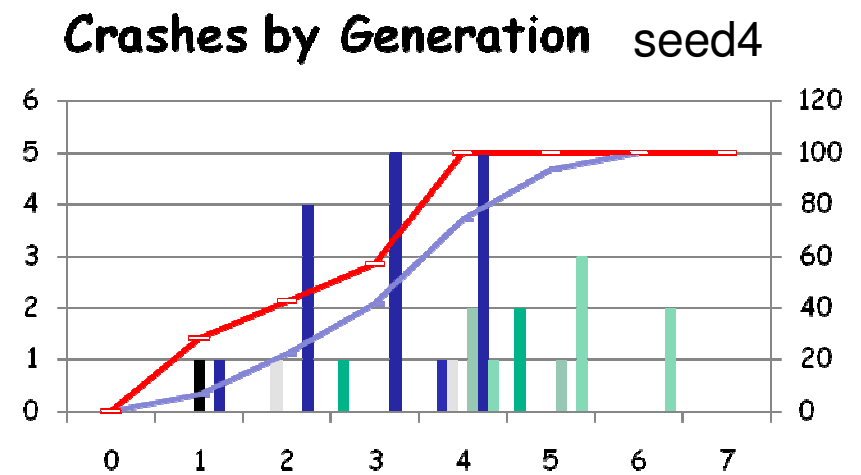
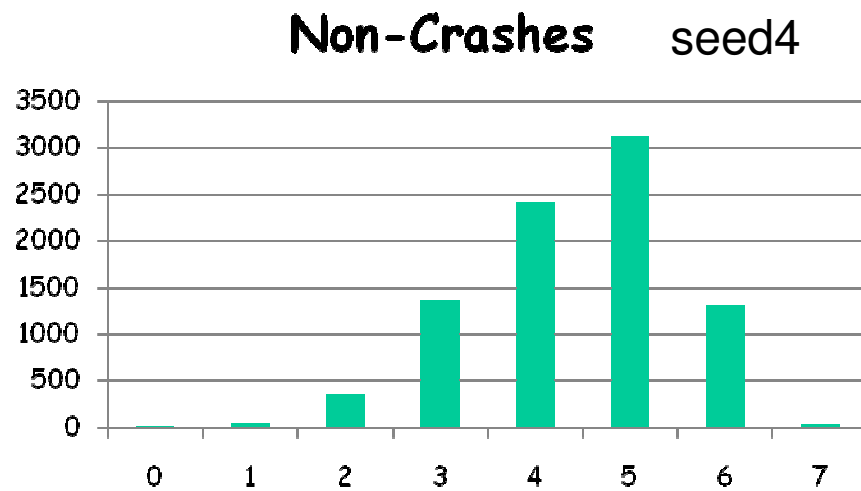
Different Seed Files, Different Crashes

| Bucket | seed1 | seed2 | seed3 | seed4 | seed5 | 100 zero bytes |
|------------|-------|-------|-------|-------|-------|----------------------|
| 1867196225 | X | X | X | X | X | |
| 2031962117 | X | X | X | X | X | |
| 612334691 | | X | X | | | |
| 1061959981 | | | X | X | | |
| 1212954973 | | | X | | | X |
| 1011628381 | | | X | X | | X |
| 842674295 | | | | X | | |
| 1246509355 | | | X | X | | X |
| 1527393075 | | | | | X | |
| 1277839407 | | | | | X | |
| 1951025690 | | | X | | | |

For the first time, we face bug triage issues!

Media1: 60 machine-hours, 44598 total tests, 357 crashes, 12 unique buckets

Most Bugs Found are "Shallow"



SAGE Summary

- SAGE is so effective at finding bugs that, **for the first time**, we face "bug triage" issues with dynamic test generation
- What makes it so effective?
 - Works on large applications (not unit test)
 - Can detect bugs due to problems across components
 - Fully automated (focus on file fuzzing)
 - Easy to deploy (x86 analysis - any language or build process !)
 - Now, used daily in various groups inside Microsoft

More On the Research Behind SAGE

- Challenges:
 - How to recover from **imprecision** in symbolic execution? PLDI'05
 - How to **scale** symbolic exec. to billions of instructions? NDSS'08
 - How to check efficiently **many properties** together? EMSOFT'08
 - How to leverage gram. specs for **complex** input formats? PLDI'08
 - How to deal with **path explosion** in large prgms? POPL'07, TACAS'08
 - How to reason **precisely** about pointers? ISSTA'09
- + research on **constraint solvers** (Z3, disolver,...)

Extension: Active Property Checking

- Traditional property checkers are “passive”
 - Purify, Valgrind, AppVerifier, TruScan, etc.
 - Check only the current concrete execution
 - Can check many properties at once
- Combine with symbolic execution “active”
 - Reason about all inputs on same path
 - Apply heavier constraint solving/proving
 - “Actively” look for input violating property
- Ex: array ref $a[i]$ where i depends on input, a is of size c
 - Try to force buffer over/underflow: add “ $(i < 0)$ OR $(i \geq c)$ ” to the path constraint; if SAT, next test should hit a bug!
- Challenge: inject/manage all such constraints efficiently...

Ext.: Grammar-Based Whitebox Fuzzing

- Input precondition specified as a context-free grammar
- Avoids path explosion in lexer and parser



- Faster, better and deeper coverage for applications with structured inputs (XML, etc.)

| generation strategy (each ran 2 hours) | #inputs | total coverage | coverage in code gen |
|---|---------|-------------------|-------------------------|
| blackbox fuzzing | 8658 | 14% | 51% |
| whitebox fuzzing | 6883 | 15% | 54% |
| grammar-based blackbox fuzzing | 7837 | 12% | 61% |
| grammar-based whitebox fuzzing | 2378 | 20% | 82% |

Ext.: Compositionality = Key to Scalability

- Problem: executing all feasible paths does not scale !
- Idea: **compositional** dynamic test generation
 - use **summaries** of individual functions (arbitrary program blocks) like in interprocedural static analysis
 - If f calls g , test g separately, summarize the results, and use g 's summary when testing f
 - A summary $\varphi(g)$ is a disjunction of path constraints expressed in terms of input preconditions and output postconditions:
$$\varphi(g) = \vee \varphi(w) \quad \text{with} \quad \varphi(w) = \text{pre}(w) \wedge \text{post}(w)$$

expressed in terms of g 's inputs and outputs
 - g 's outputs are treated as symbolic inputs to a calling function f
- Can provide same path coverage exponentially faster !

Conclusion: Blackbox vs. Whitebox Fuzzing

- Different cost/precision tradeoffs
 - Blackbox is lightweight, easy and fast, but poor coverage
 - Whitebox is smarter, but complex and slower
 - Note: other recent "semi-whitebox" approaches
 - Less smart (no symbolic exec, constr. solving) but more lightweight: Flayer (taint-flow, may generate false alarms), Bunny-the-fuzzer (taint-flow, source-based, fuzz heuristics from input usage), etc.
- Which is more effective at finding bugs? It depends...
 - Many apps are so buggy, any form of fuzzing find bugs in those !
 - Once low-hanging bugs are gone, fuzzing must become smarter: use whitebox and/or user-provided guidance (grammars, etc.)
- Bottom-line: in practice, use both! (We do at Microsoft)

Future Work (The Big Picture)

- During the last decade, **code inspection** for **standard programming errors** has largely been **automated** with **static code analysis**
- Next: **automate testing** (as much as possible)
 - Thanks to advances in program analysis, efficient constraint solvers and powerful computers
- Whitebox testing: automatic code-based test generation
 - Like static analysis: automatic, scalable, checks many properties
 - Today, we can exhaustively test small applications, or partially test large applications
 - Biggest impact so far: whitebox fuzzing for (Windows) security testing
 - Improved security for a billion computers worldwide!
 - Next: towards exhaustive testing of large application (**verification**)
 - How far can we go?